

Parsing Challenges in Java 8

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JastAddJ is a full source-to-bytecode modular Java compiler

- each Java version is a separate module
- Java 8 was implemented by Erik Hogeman for his Master's Thesis
- this talk is about the parsing challenges encountered

Noteworthy features:

- Lambdas
- Method references
- Default methods

Java finally has anonymous functions!

$(x, y) \rightarrow x + y$

$() \rightarrow \{ \text{action1}(); \text{action2}(); \}$

Lambda Example

Action listeners the old way:

```
button.addActionListener(new ActionListener() {  
    public void actionPerformed(ActionEvent e) {  
        print("hello");  
    }  
});
```

The new way, using lambda:

```
button.addActionListener( (e) -> print("hello") );
```



A way of using regular instance methods as lambdas:

```
Greeter greeter = new MyGreeter();  
greetButton.addActionListener( greeter::greet );  
exitButton.addActionListener( greeter::exit );
```

Default Methods

Interfaces can have non-abstract methods:

```
interface Greeter {  
    default void greet(ActionEvent e) {  
        print("greetings");  
    }  
    default void exit(ActionEvent e) {  
        print("goodbye");  
    }  
}  
  
class MyGreeter implements Greeter {  
    // use default implementations  
}
```



We use an LALR parser for JastAddJ

- Generated with the Beaver parser generator
- Parser grammar is composed from parts in separate modules

Why an LR Parser Generator?

Advantages of a generated LR parser:

- Provably fast
- Generator certifies unambiguous grammar
- Decent tool support
- Bit more powerful than LL



Java 8 Parsing Challenges

- Ambiguous grammar specification
- Reduce-reduce conflicts between subexpressions
- Shift-reduce conflict
- Unlimited lookahead



Ambiguous Grammar Specification

Java spec (highly edited):

Expression \rightarrow Lambda

Expression \rightarrow ... \rightarrow Additive

\rightarrow Multiplicative \rightarrow ... \rightarrow Cast

Cast \rightarrow (Type) Lambda

Input:

(T) (a, b) \rightarrow a * b;

Possible parse 1:

((T) (a, b) \rightarrow a) * b;

Possible parse 2:

(T) ((a, b) \rightarrow a * b);

The second one is desired. We achieved this by:

- changed the grammar
- lambda as primary expression
- lowered priority using precedence declarations



Lambda Reduce-Reduce Conflict

Lambda vs less-than expression:

```
(T<A> s) -> { } // lambda  
(T<A)           // less-than expression
```

This is a reduce-reduce conflict.

Similar conflict in Java 5 with type cast:

```
(T<A>) s           // generic type cast  
(T<A)             // less-than expression
```

In both cases the T terminal must be reduced to either *RelationalExpression* or *ReferenceType*.



Lambda Reduce-Reduce Conflict

We solved the reduce-reduce conflict by giving the related parsing productions explicit common prefixes:

Relational -> Name < Shift

Relational -> Relational < Shift

...

ReferenceType -> Name < TypeArguments_1

This removed the need to reduce the Name token too early.

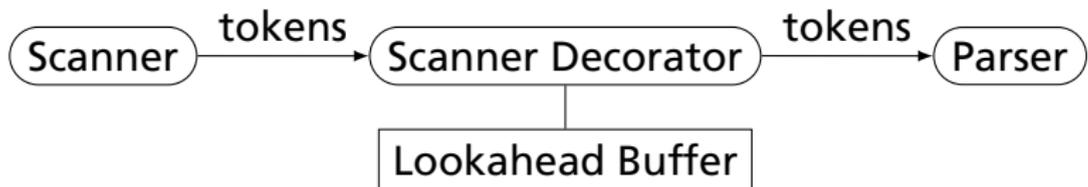


Unlimited Lookahead

```
f(T<A, B>::m)    // method reference  
f(T<A, B> m)     // less-than expression
```

There is no reasonable fixed lookahead that will allow the parser to decide between a less-than expression, or method reference.

Scanner Decorator



The Scanner Decorator looks ahead in the token stream when certain tokens are encountered, then potentially modifies the token stream.

In the previous case it inserts a synthetic `LT_TYPE` token.

Conclusions

- Java is not LR, but with some modifications we can make it LR(1)
- So far implemented nearly all of Java 8 features (parsing is complete)

Techniques we used to solve parsing challenges:

- Duplicate grammar to avoid reduce-reduce conflicts
- Introduce priority declarations to fix ambiguous grammar
- Scanner decorator to enable infinite lookahead



Questions!

Default Modifier Shift-Reduce

We parse all modifiers using the same production (for methods, interfaces, classes).

This introduced a shift-reduce conflict in switch-statements:

```
switch (x) {  
  case 0:  
    default class A() { };  
  case 1:  
    break;  
  default:  
}
```

Intersection Type Cast

In Java 8 cast expressions can have the form:

`(A & B & C) x`

This form conflicts with binary expressions:

`(A & B & C)`

The conflict is very similar to the lambda versus less-than expression conflict.

Parsing Intersection Type Casts

We solve this conflict using the Scanner Decorator. Whenever a left-parenthesis is encountered, the decorator inserts the synthetic INTERCAST token if it determines that it is part of an intersection type cast.